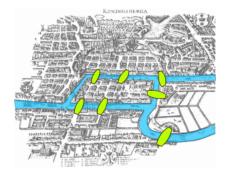
Graph Theory II

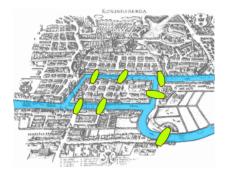
Nashville Math Club

October 20, 2020

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Question

Can we walk through all seven bridges of Konigsberg exactly once?

Lemma

The sum of the degrees of all the vertices in a graph is equal to two times the number of edges in the graph.

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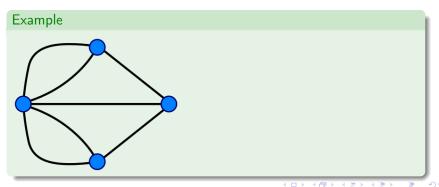
This means there must always be an even number of odd vertices!

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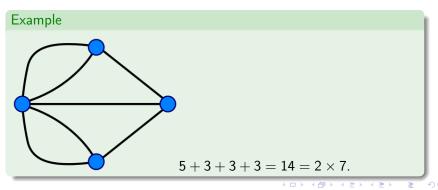
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Each landmass is a vertex and each bridge is an edge.

Question

Can we walk through all seven bridges of Konigsberg exactly once?



Each landmass is a vertex and each bridge is an edge. Notice that whenever one enters a vertex by an edge they must leave the vertex by a different edge.

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Can we walk through all seven bridges of Konigsberg exactly once?



Each landmass is a vertex and each bridge is an edge. Notice that whenever one enters a vertex by an edge they must leave the vertex by a different edge.

Therefore, each vertex (except for the starting and finishing ones) must have an even degree.

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Can we walk through all seven bridges of Konigsberg exactly once?



Each landmass is a vertex and each bridge is an edge. Notice that whenever one enters a vertex by an edge they must leave the vertex by a different edge.

Therefore, each vertex (except for the starting and finishing ones) must have an even degree.

The walk is only possible if exactly zero or two vertices have odd degree.

Definition

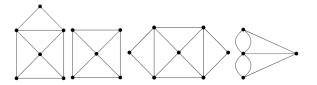
• An **Euler path** in a graph or multigraph is a path which uses each edge exactly once.

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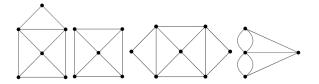
Definition

- An **Euler path** in a graph or multigraph is a path which uses each edge exactly once.
- An **Euler circuit** is an Euler path which starts and stops at the same vertex.

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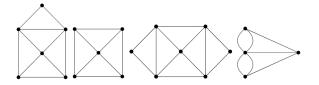


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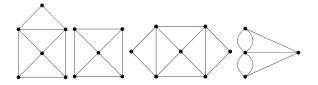
Question

• Which of these graphs have Euler paths or circuits?



Question

- Which of these graphs have Euler paths or circuits?
- Is there a connection between the degrees of each vertex and your answer above?



Question

- Which of these graphs have Euler paths or circuits?
- Is there a connection between the degrees of each vertex and your answer above?
- How would you draw a graph that's guaranteed to have an Euler circuit?

Theorem

• A connected graph has an Euler circuit if and only if the degree of each vertex is even.

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Theorem

- A connected graph has an Euler circuit if and only if the degree of each vertex is even.
- A connected graph has an Euler path if and only if there are at most two vertices with odd degree.

Suppose instead of traveling on each bridge in Konigsberg you wanted instead to travel to each landmass exactly once.

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Suppose instead of traveling on each bridge in Konigsberg you wanted instead to travel to each landmass exactly once.

Definition

• A Hamilton path is a path which visits each vertex of a graph exactly once.

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Suppose instead of traveling on each bridge in Konigsberg you wanted instead to travel to each landmass exactly once.

Definition

- A **Hamilton path** is a path which visits each vertex of a graph exactly once.
- A Hamilton cycle is a Hamilton path which starts and stops at the same vertex.

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Question

Is there a Hamilton path for Konigsberg?



Question

Is there a Hamilton path for Konigsberg?

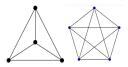
Question

How can we determine if there is a Hamilton path for a general graph G?

Recall K_n is the graph with n vertices where each pair of distinct vertices is connected by a unique edge.

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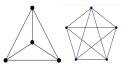
Recall K_n is the graph with n vertices where each pair of distinct vertices is connected by a unique edge.



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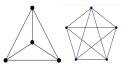
Question

• For which n does K_n have an Euler path? An Euler circuit?

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Recall K_n is the graph with n vertices where each pair of distinct vertices is connected by a unique edge.



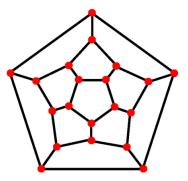
Question

- For which n does K_n have an Euler path? An Euler circuit?
- For which n does K_n have a Hamilton path? A Hamilton cycle?

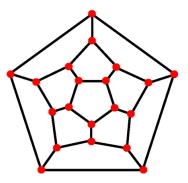
Recall that a graph is **planar** if its edges only intersect at vertices and is **connected** if theres a path from each vertex to every other vertex.

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Let v be the number of vertices, e the number of edges, and f the number of faces or regions.

Draw several planar connected graphs and compute v, e, and f for each of them.

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Draw several planar connected graphs and compute v, e, and f for each of them. Do you notice any patterns?

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Draw several planar connected graphs and compute v, e, and f for each of them.

Do you notice any patterns?

Definition

The Euler Characteristic χ is defined by

$$\chi := \mathbf{v} - \mathbf{e} + \mathbf{f}.$$

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Question

What is the Euler characteristic of the graphs you drew?

Recall a tree is a graph with no cycles.

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Recall a tree is a graph with no cycles.

Question

What is the Euler characteristic of a tree?

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Recall a tree is a graph with no cycles.

Question

What is the Euler characteristic of a tree?

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Notice that all trees have v = e + 1.

Recall a tree is a graph with no cycles.

Question

What is the Euler characteristic of a tree?

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Notice that all trees have v = e + 1. What is f for a tree?

Recall a tree is a graph with no cycles.

Question

What is the Euler characteristic of a tree?

Notice that all trees have v = e + 1. What is f for a tree? f = 1 for all trees.

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Recall a tree is a graph with no cycles.

Question

What is the Euler characteristic of a tree?

Notice that all trees have v = e + 1. What is f for a tree? f = 1 for all trees. Therefore all trees have Euler characteristic

$$\chi = \mathbf{v} - \mathbf{e} + \mathbf{f} = 2.$$

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Theorem

The Euler characteristic of every planar connected graph is 2.

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Theorem

The Euler characteristic of every planar connected graph is 2.

Proof

 If you have a connected planar graph that isn't a tree, remove an edge that completes a cycle.

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Theorem

The Euler characteristic of every planar connected graph is 2.

Proof

• If you have a connected planar graph that isn't a tree, remove an edge that completes a cycle.

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• What does this do to the Euler characteristic?

Theorem

The Euler characteristic of every planar connected graph is 2.

Proof

• If you have a connected planar graph that isn't a tree, remove an edge that completes a cycle.

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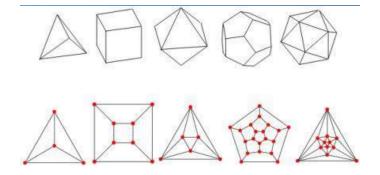
- What does this do to the Euler characteristic?
- We can do this until we are left with a tree.

Theorem

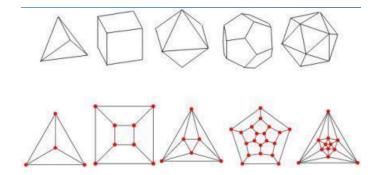
The Euler characteristic of every planar connected graph is 2.

Proof

- If you have a connected planar graph that isn't a tree, remove an edge that completes a cycle.
- What does this do to the Euler characteristic?
- We can do this until we are left with a tree.
- Therefore the Euler characteristic of all connected planar graphs is 2.

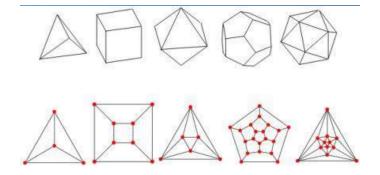


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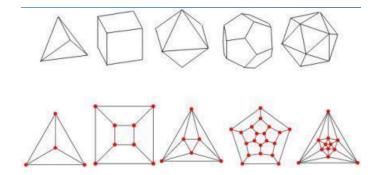


Each platonic solid has a corresponding planar connected graph

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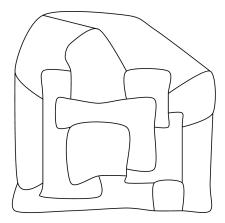


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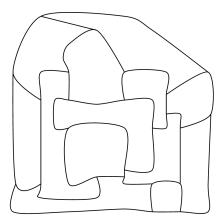


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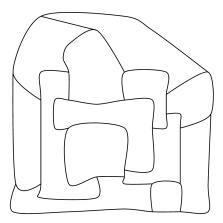


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Mapmakers want to color this map, but want each adjacent region to have a different color.

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Mapmakers want to color this map, but want each adjacent region to have a different color. How many colors are needed?

Definition

• Given a graph G, a coloring of the vertices is called a vertex coloring

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Definition

- Given a graph *G*, a coloring of the vertices is called a **vertex coloring**
- If all adjacent vertices are colored differently, the coloring is called **proper**.

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Question

Does every graph have a proper coloring?

Definition

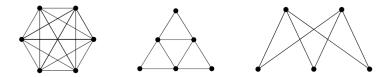
- Given a graph G, a coloring of the vertices is called a vertex coloring
- If all adjacent vertices are colored differently, the coloring is called **proper**.

Question

Does every graph have a proper coloring?

Definition

The smallest number of colors needed to get a proper vertex coloring is called the **chromatic number** of the graph.

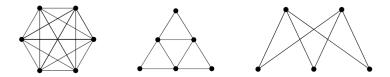


Question

• What are the chromatic numbers of each of these graphs?

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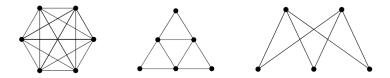
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Question

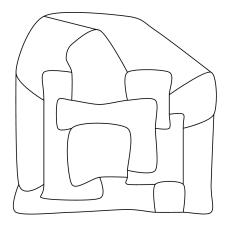
- What are the chromatic numbers of each of these graphs?
- For each positive integer n, is there a graph with chromatic number equal to n?

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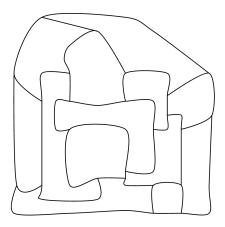


Question

- What are the chromatic numbers of each of these graphs?
- For each positive integer n, is there a graph with chromatic number equal to n?
- What is the chromatic number of a tree?



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If we represent each region by a vertex and adjacent regions are connected by an edge what properties does our graph have?

Theorem (The four color theorem)

If G is a planar graph, then its chromatic number is less than or equal to 4.

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Theorem (The four color theorem)

If G is a planar graph, then its chromatic number is less than or equal to 4.

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Therefore every map can be colored with 4 or less colors!

Question

What is the smallest number of colors that can be used to color the vertices of a cube so that no adjacent vertices have the same color?

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