CS 4260 and CS 5260 Vanderbilt University

Lecture on Adversarial Search in Games

This lecture covers material in optional reading from section 11.3 of ArtInt

ArtInt: Poole and Mackworth, Artificial Intelligence 2E at http://artint.info/2e/html/ArtInt2e.html to include slides at http://artint.info/2e/slides/ch04/lect1.pdf

In game of Othello black moves first, but where?



Search ahead

Game ends when no player has a move





Using simple difference utility function doesn't discriminate at shallow levels

Some squares are more important than others

Use a difference of weighted sums

Due to Peter Norvig in "Artificial Intelligence Programming" 1992, Morgan Kaufmann Publishers





Probably not! Opponent won't cooperate





Assumptions: larger scores better for player (Max). smaller scores better for opponent (Min) utility function is rational



Minimax search with no pruning



Minimax search with no pruning



Minimax search with no pruning

All this reasoning is going on in "mind" of AI about to move





Is there any need to examine this node (i.e., game board)?

No – the maximizing player will not move in that direction anyways



Minimax search with (apha beta) pruning















Take middle move



This slide is specific to sample Othello code – you won' be examined on it! This is a C implementation of a variation on a Lisp version by Peter Norvig in "Artificial Intelligence Programming" 1992, Morgan Kaufmann Publishers



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Finally, remember that in Othello, player **i** may not have any legal moves, so the other player, **k**, continues to move until player **i** can once again move. As a result, *minchoice* will sometimes call *minchoice* and *maxchoice* will sometimes call *maxchoice*.



If you run roundrobin with the current code (minmax with no pruning), then be prepared to wait quite a while. The following will eventually materialize (which I have indented for a bit better readability):

random wins=1 boards=0 random wins=2 boards=0 random wins=0 boards=0 random wins=3 boards=0 random wins=1 boards=0 random wins=0 boards=0 diff/1 wins=0 boards=1799 diff/1 wins=0 boards=1763 diff/1 wins=4 boards=2134 diff/1 wins=2 boards=2038 diff/1 wins=0 boards=1736 diff/3 wins=3 boards=200450 diff/3 wins=3 boards=202781 diff/3 wins=1 boards=208165 diff/3 wins=1 boards=174285 diff/5 wins=2 boards=18595500 diff/5 wins=4 boards=19102429 diff/5 wins=0 boards=16349918 wdiff/1 wins=0 boards=2311 wdiff/1 wins=1 boards=1944 wdiff/3 wins=1 boards=234330

diff/1 wins=9 boards=1824 diff/3 wins=8 boards=174938 diff/5 wins=10 boards=14970716 wdiff/1 wins=7 boards=2455 wdiff/3 wins=9 boards=293989 wdiff/5 wins=10 boards=33035881 diff/3 wins=10 boards=155657 diff/5 wins=10 boards=16361232 wdiff/1 wins=6 boards=2397 wdiff/3 wins=7 boards=268585 wdiff/5 wins=10 boards=24682590 diff/5 wins=6 boards=19828063 wdiff/1 wins=7 boards=2301 wdiff/3 wins=9 boards=242318 wdiff/5 wins=9 boards=26257282 wdiff/1 wins=8 boards=2383 wdiff/3 wins=5 boards=202807 wdiff/5 wins=10 boards=22520544 wdiff/3 wins=10 boards=292058 wdiff/5 wins=9 boards=29422051 wdiff/5 wins=8 boards=28452315

Each line above shows the results of 10 matches against two strategies. For example, the last line shows the results of the strategy using 3-ply lookahead with weighteddiffeval (wdiff/3) as the utility function against 5-ply lookahead with the same utility function (wdiff/5). Over 10 games, wdiff/5 wins 8, wdiff/3 wins 1, and there is one tie. Over 10 games, wdiff/5 (in this last case) generates 28,452,315 boards (or an average of 2,845,232 boards generated per game)!

```
diff1 wins=9 boards=1824
random wins=1 boards=0
                               diff3 wins=8 boards=44887
random wins=2 boards=0
                               diff5 wins=10 boards=944422
random wins=0 boards=0
random wins=3 boards=0
                               wdiff1 wins=7 boards=2455
random wins=1 boards=0
                               wdiff3 wins=9 boards=97418
random wins=0 boards=0
                               wdiff5 wins=10 boards=2839502
                               diff3 wins=10 boards=44741
diff1 wins=0 boards=1799
diff1 wins=0 boards=1763
                               diff5 wins=10 boards=1038063
diff1 wins=4 boards=2134
                               wdiff1 wins=6 boards=2397
diff1 wins=2 boards=2038
                               wdiff3 wins=7 boards=90605
diff1 wins=0 boards=1736
                               wdiff5 wins=10 boards=2643711
                               diff5 wins=6 boards=1160722
diff3 wins=3 boards=59175
diff3 wins=3 boards=51624
                               wdiff1 wins=7 boards=2301
                               wdiff3 wins=9 boards=84970
diff3 wins=1 boards=58869
diff3 wins=1 boards=48521
                               wdiff5 wins=9 boards=2947041
diff5 wins=2 boards=1142881
                               wdiff1 wins=8 boards=2383
diff5 wins=4 boards=1063738
                               wdiff3 wins=5 boards=76715
diff5 wins=0 boards=1124236
                               wdiff5 wins=10 boards=2683518
wdiff1 wins=0 boards=2311
                               wdiff3 wins=10 boards=89526
wdiff1 wins=1 boards=1944
                               wdiff5 wins=9 boards=2635355
wdiff3 wins=1 boards=80708
                               wdiff5 wins=8 boards=2844060
```

Notice that the win/loss results do not change at all (minmax with alphabeta pruning chooses exactly the same moves as minmax without alphabeta pruning), but the number of generated boards is vastly reduced. For example, when wdiff5 is pitted against wdiff3, wdiff5 with pruning generates 2,844,060 boards over 10 games or an average of 284,406 boards per game, as opposed to about 10 times this without pruning.

Exercise: understand the book's algorithm on examples

1:pr	ocedure Minimax_alpha_beta(n, α, β)	
2:	Inputs	
3:	n a node in a game tree	Minimar Alpha Beta $(R - \infty \infty)$
4:	α, β real numbers	$Minimuz Alphu Deta (n, -\infty, \infty)$
5:	Output	
6:	A pair of a value for node n, path that gives this value	Max
7:	best := None	
8:	if n is a leaf node then	
9:	return $evaluate(n)$, None	
10:	else if n is a MAX node .hen	
11:	for each child c of n do	
12:	score, path := MinimaxAlphaBeta(c, lpha, eta)	
13:	if $score \geq \beta$ then	
14:	return score, None	
15:	else if $score > \alpha$ then	
16:	$\alpha := score$	
17:	best := c : path	
18:	return α , best	
19:	else	
20:	for each child c of n do	
21:	$score, path := MinimaxAlphaBeta(c, \alpha, \beta)$	$2 \qquad 7 \qquad 1 \qquad 0$
22:	if $score \leq \alpha$ then	\angle 1 1 9
23:	return score, None	
24:	else if $score < \beta$ then	
25:	$\beta := score$	
26:	best := c : path	
27:	return β , best	
	Figure 11.6: Minimax with $\alpha - \beta$ pruning	

From ArtInt (Poole and Mackworth) Section 11.3





1:pr	ocedure Minimax_alpha_beta(n, α, β)			
2:	Inputs			
3:	n a node in a game tree			· · (D)
4:	lpha,eta real numbers		MinimaxAlphaB	$eta(R,-\infty,\infty)$
5:	Output			
6:	A pair of a value for node <i>n</i> , path that gives	this value		Max
7:	best := None			
8:	if n is a leaf node then			<
9:	return $evaluate(n)$, None	Minim	$ar Alpha Beta (B - \infty \infty)$	\mathbf{i}
10:	else if n is a MAX node then	111 0100110	$(10, \infty, \infty)$	
11:	for each child c of n do			\searrow
12:	score, path := MinimaxAlphaBeta	(c, α, β)		Min
13:	if $score \geq \beta$ then			
14:	return score, None			
15:	else if $score > \alpha$ then		2 ′ / \	
16:	$\alpha := score$			
17:	best := c : path Mini	maxAlphal	$Beta(R, -\infty, \infty)$	
18:	return α , best	-		
19:	else			\sim
20:	for each child c of n do		Max	\bigcirc
21:	score, path := MinimaxAlphaBeta	(c, α, β)		1 0
22:	if $score \leq \alpha$ then		Z 1	1 9
23:	return score, None			
24:	else if $score < \beta$ then			
25:	$\beta := score$			
26:	best := c : path			
27:	return β , best			
	Figure 11.6: Minimax with α - β pruning	ng		



1:pr	ocedure Minimax_alpha_beta(n, α, β)		
2:	Inputs		
3:	n a node in a game tree		
4:	lpha,eta real numbers		$MinimaxAlphaBeta\left(R,-\infty,\infty ight)$
5:	Output		
6:	A pair of a value for node n, path that gives	this value	Max
7:	best := None		
8:	if n is a leaf node then		
9:	return $evaluate(n)$, None	Minim	$ar Alpha Beta (B - \infty \alpha)$
10:	else if n is a MAX node then	141 0100110	$(1, -\infty, \infty)$
11:	for each child c of n do		\sim 2
12:	score, path := MinimaxAlphaBeta	(c, α, β)	\bigcirc \checkmark \checkmark \bigcirc Min
13:	if $score \geq \beta$ then		
14:	return score, None		
15:	else if $score > \alpha$ then		
16:	$\alpha := score$		
17:	best := c : path		$/$ Minimax Alpha Beta $(R, -\infty, 2)$
18:	return α , best		
19:	else		
20:	for each child c of n do		\bigcirc Max \bigcirc \bigcirc
21:	score, path := MinimaxAlphaBeta	(c, α, β)	$2 \qquad 7 \qquad 1 \qquad 0$
22:	if $score \leq \alpha$ then		Z / I 9
23:	return score, None		
24:	else if $score < \beta$ then		
25:	$\beta := score$		
26:	best := c : path		
27:	return β , best		
	Figure 11.6: Minimax with α - β pruni	ng	

1:pr	ocedure Minimax_alpha_beta(n, α, β)		
2:	Inputs		
3:	n a node in a game tree		
4:	α, β real numbers		$MinimaxAlphaBeta\left(R,-\infty,\infty ight)$
5:	Output		
6:	A pair of a value for node n , path that gives	this value	Max
7:	best := None		
8:	if n is a leaf node then		
9:	return $evaluate(n)$, None	Minim	$ar Alpha Beta (B - \infty , \infty)$
10:	else if n is a MAX node then	111 111111	$(n, -\infty, \infty)$
11:	for each child c of n do		
12:	score, path := MinimaxAlphaBeta	(c, α, β)	
13:	if $score \geq \beta$ then		
14:	return score, None		
15:	else if $score > \alpha$ then		
16:	$\alpha := score$		
17:	best := c : path		$Minimar Alpha Beta (B - \infty, 2)$
18:	return α , best		
19:	else		
20:	for each child c of n do		\bigcirc Max \bigcirc \bigcirc
21:	score, path := MinimaxAlphaBeta	(c, α, β)	2 7 4 0
22:	if $score \leq \alpha$ then		2 / 1 9
23:	return score, None		
24:	else if $score < \beta$ then		$7 \neq \infty$
25:	$\beta := score$		/ 7 -00
26:	best := c : path		7 4 2
27:	return β , best		
	Figure 11.6: Minimax with α - β pruni	ng	

1:pr	ocedure Minimax_alpha_beta(n, α, β)			
2:	Inputs			
3:	n a node in a game tree		16: : 41.1	D (/ D)
4:	lpha,eta real numbers		MinimaxAlpho	$iBeta\left(R,-\infty,\infty ight)$
5:	Output			
6:	A pair of a value for node <i>n</i> , path that gives	this value		Max
7:	best := None		2	
8:	if n is a leaf node then			
9:	return $evaluate(n)$, None	Minin	par Alpha Beta R - ~ d	
10:	else if n is a MAX node then	141 010011	counterinancea (n.,, A	"
11:	for each child c of n do		\sim 2	
12:	score, path := MinimaxAlphaBeta	(c, α, β)		Min
13:	if $score \ge \beta$ then			
14:	return score, None			
15:	else if $score > \alpha$ then			
16:	$\alpha := score$			
17:	best := c : path			
18:	return $\alpha, best$			
19:	else // n is a MIN node			
20:	for each child c of n do		Max	
21:	score, path := MinimaxAlphaBeta	(c, α, β)	\mathbf{O} $\mathbf{\nabla}$	1 0
22:	if $score \leq \alpha$ then		Δ /	1 9
23:	return score, None			
24:	else if $score < \beta$ then			
25:	$\beta := score$			
26:	best := c : path			
27:	return $\beta, best$			
	Figure 11.6: Minimax with α - β prunimation	ng		

1:pr	ocedure Minimax_alpha_beta(n, α, β)	
2:	Inputs	
3:	n a node in a game tree	
4:	α, β real numbers	$Minimax Alpha Beta (R, -\infty)$
5:	Output	112 (100 / 102 (10, 1 0, 00)
6:	A pair of a value for node n, path that gives this value	\sim \sim 2
7:	best := None	Max 2
8:	if n is a leaf node then	
9:	return evaluate(n), None	
10:	else if n is a MAX node then	$MinimaxAlphaBeta(R,\ 2\ ,\ \infty)$
11:	for each child c of n do	
12:	score, path := MinimaxAlphaBeta(c, lpha, eta)	
13:	if $score \geq \beta$ then	Min View Min
14:	return score, None	
15:	else if $score > \alpha$ then	
16:	$\alpha := score$	
17:	best := c : path	
18:	return α , best	
19:	else	
20:	for each child c of n do	
21:	$score, path := MinimaxAlphaBeta(c, \alpha, \beta)$	
22:	if $score \leq \alpha$ then	2 7 1 9
23:	return score, None	
24:	else if $score < \beta$ then	
25:	$\beta := score$	
26:	best := c : path	
27:	return β , best	
	Figure 11.6: Minimax with $\alpha - \beta$ pruning	

1:pr	ocedure Minimax_alpha_beta(n, α, β)	
2:	Inputs	
3:	n a node in a game tree	
4:	α, β real numbers	$Minimax Alpha Beta (R, -\infty, \infty)$
5:	Output	
6:	A pair of a value for node n , path that gives this value	\sim \sim 2
7:	best := None	Max 2
8:	if n is a leaf node then	
9:	return $evaluate(n)$, None	
10:	else if n is a MAX node then	$MinimaxAlphaBeta(R,\ 2\ ,\ \infty)$
11:	for each child c of n do	
12:	score, path := MinimaxAlphaBeta(c, lpha, eta)	
13:	if $score \geq \beta$ then	Min
14:	return score, None	
15:	else if $score > \alpha$ then	
16:	$\alpha := score$	
17:	best := c : path	
18:	return α , best	
19:	else	
20:	for each child c of n do	
21:	$score, path := MinimaxAlphaBeta(c, \alpha, \beta)$	
22:	if $score \leq \alpha$ then	2 7 1 9
23:	return score, None	
24:	else if $score < \beta$ then	
25:	$\beta := score$	
26:	best := c : path	
27:	return β , best	
	Figure 11.6: Minimax with α - β pruning	

1:pr	ocedure Minimax_alpha_beta(n, α, β)	
2:	Inputs	
3:	n a node in a game tree	
4:	lpha,eta real numbers	$Minimax Alpha Beta (R, -\infty, \infty)$
5:	Output	
6:	A pair of a value for node <i>n</i> , path that gives this value	\sim \sim 2
7:	best := None	Max ²
8:	if n is a leaf node then	
9:	return evaluate(n), None	
10:	else if n is a MAX node then	$MinimaxAlphaBeta(R,\ 2\ ,\ \infty)$
11:	for each child c of n do	
12:	$score, path := MinimaxAlphaBeta\left(c, lpha, eta ight)$	
13:	if $score \ge \beta$ then	Min
14:	return score, None	
15:	else if $score > \alpha$ then	
16:	$\alpha := score$	
17:	best := c : path	\wedge MinimaxAlphaBeta (R, 2, ∞)
18:	return α , best	
19:	else	
20:	for each child c of n do	
21:	score, path := MinimaxAlphaBeta(c, lpha, eta)	
22:	if $score \leq \alpha$ then	2 7 1 9
23:	return score, None	
24:	else if $score < \beta$ then	
25:	$\beta := score$	
26:	best := c : path	
27:	return β , best	
	Figure 11.6: Minimax with α - β pruning	

1:pro	ocedure Minimax_alpha_beta(n, α, β)	
2:	Inputs	
3:	n a node in a game tree	
4:	lpha,eta real numbers	$Minimax Alpha Beta (R, -\infty, \infty)$
5:	Output	
6:	A pair of a value for node <i>n</i> , path that gives this value	\sim \sim 2
7:	best := None	Max 2
8:	if n is a leaf node then	
9:	return evaluate (n), None	
10:	else if n is a MAX node then	$MinimaxAlphaBeta(R, 2, \infty)$
11:	for each child c of n do	
12:	score, path := MinimaxAlphaBeta(c, lpha, eta)	
13:	if $score \geq \beta$ then	
14:	return score, None	
15:	else if $score > \alpha$ then	
16:	$\alpha := score$	
17:	best := c : path	\wedge MinimaxAlphaBeta (R, 2, ∞)
18:	return α , best	
19:	else	
20:	for each child c of n do	Max C
21:	$score, path := MinimaxAlphaBeta(c, \alpha, \beta)$	
22:	if $score \leq \alpha$ then	2 7 1 9
23:	return score, None	
24:	else if $score < \beta$ then	1 < 2
25:	$\beta := score$	$1 \leq 2$
26:	best := c : path	
27:	return β , best	
	Figure 11.6: Minimax with α - β pruning	



1:pr	ocedure <i>Minimax_alpha_beta</i> (n, α, β)	
2:	Inputs	
3:	n a node in a game tree	
4:	α, β real numbers	$Minimar Alpha Beta (B, -\infty)$
5:	Output	
6:	A pair of a value for node n, path that gives this value	\sim \sim 2
7:	best := None	Max ²
8:	if n is a leaf node then	$1 \neq \infty$
9:	return $evaluate(n)$, None	
10:	else if n is a MAX node then	$1 \neq 2$
11:	for each child c of n do	
12:	$score, path := MinimaxAlphaBeta\left(c, lpha, eta ight)$	
13:	if $score \geq \beta$ then	Min Min
14:	return score, None	
15:	else if $score > \alpha$ then	
16:	$\alpha := score$	
17:	best := c : path	
18:	return α , best	
19:	else	
20:	for each child c of n do	Max Max
21:	$score, path := MinimaxAlphaBeta(c, \alpha, \beta)$	
22:	if $score \leq \alpha$ then	2 7 1 9
23:	return score, None	
24:	else if $score < \beta$ then	
25:	$\beta := score$	
26:	best := c : path	
27:	return β , best	
	Figure 11.6: Minimax with $\alpha - \beta$ pruning	

1:pr	ocedure Minimax_alpha_beta(n, α, β)	2
2:	Inputs	
3:	n a node in a game tree	
4:	α, β real numbers	$Minimar Alpha Beta (B - \infty)$
5:	Output	$(10, \infty, \infty)$
6:	A pair of a value for node <i>n</i> , path that gives this value	1 2
7:	best := None	\bigcirc Max \checkmark
8:	if n is a leaf node then	
9:	return $evaluate(n)$, None	
10:	else if n is a MAX node then	
11:	for each child c of n do	
12:	score, path := $MinimaxAlphaBeta(c, \alpha, \beta)$	
13:	if $score \geq \beta$ then	Min
14:	return score, None	
15:	else if $score > \alpha$ then	
16:	$\alpha := score$	
17:	best := c : path	
18:	return α , best	
19:	else	
20:	for each child c of n do	
21:	$score, path := MinimaxAlphaBeta(c, \alpha, \beta)$	
22:	if $score \leq \alpha$ then	2 7 1 9
23:	return score, None	
24:	else if $score < \beta$ then	
25:	$\beta := score$	
26:	best := c : path	
27:	return β , best	
	Figure 11.6: Minimax with $\alpha - \beta$ pruning	

9

Game Tree Search Problems

Given the following game tree, compute the utility value of the root (which is at a maximizing level) as computed by minimax search, and write the utility value by the root



Game Tree Search Problems

Put an 'X' through an arc if it and its entire subtree would be pruned when using alpha beta pruning.







Consider the game tree below. The numbers at leaves are scores obtained by an evaluation function. Show the value of the root obtained through minimax search. Additionally, put an 'X' through an arc if it and its entire subtree would be pruned when using alpha beta pruning.



Consider the game tree below. The numbers at leaves are scores obtained by an evaluation function. Show the value of the root obtained through minimax search. Additionally, put an 'X' through an arc if it and its entire subtree would be pruned when using alpha beta pruning.

