# Efficient Boltzmann Samplers for Weighted Partitions and Selections 

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## Weighted Partitions

Def. An integer partition of $n$ is a decomposition of $n$ into a nonincreasing sequence of positive integers that sums to $n$.

- Partitions of 4: (4), $(3,1),(2,2),(2,1,1),(1,1,1,1)$

Def. A weighted partition of $n$ is a partition of $n$ where each summand of size $k$ belongs to one of $b_{k}$ types. The class $\mathcal{C}$ of weighted partitions has the generating function

$$
\begin{aligned}
C(z) \equiv \sum_{n=0}^{\infty} c_{n} z^{n} & \equiv \prod_{k=1}^{\infty}\left(1-z^{k}\right)^{-b_{k}} \\
& =\prod_{k=1}^{\infty}\left(1+z^{k}+z^{2 k}+\ldots\right)^{b_{k}}
\end{aligned}
$$

- If $\left(b_{k}\right)_{k=1}^{\infty}=(1,1,3,1,1, \ldots)$ the weighted partitions of 4 are: $(4),(3,1),(3,1),(3,1),(2,2),(2,1,1),(1,1,1,1)$.


## Weighted Partitions

Example. Bose-Einstein condensation occurs when $b_{k}=\binom{k+2}{2}$.

- $C(z)=\prod_{k=1}^{\infty}\left(1-z^{k}\right)^{-\binom{k+2}{2}}=1+3 z+12 z^{2}+38 z^{3}+\ldots$

| |ll!
(3)

$(1,1,1)$
Figure: Young diagrams for $n=3$.


## Weighted Partitions

Problem. Generate a weighted partition of $n$ uniformly at random.


Figure: Random integer partition and weighted partition for $n=10^{5}$.

## Boltzmann Samplers

Def. A Boltzmann distribution over the class $\mathcal{C}$ parameterized by $0<\lambda<\rho_{\mathcal{C}}$ is the probability distribution, for all $\gamma \in \mathcal{C}$, defined as

$$
\mathbb{P}_{\lambda}(\gamma)=\frac{\lambda^{|\gamma|}}{C(\lambda)}
$$

where $\rho_{C}$ is the radius of convergence of $C(z)=\sum_{n=0}^{\infty} c_{n} z^{n}$.

Def. A Boltzmann sampler $\Gamma C(\lambda)$ is an algorithm that generates objects from $\mathcal{C}$ according to the Boltzmann distribution parameterized by $\lambda$.

- All objects of size $n$ occur with equal probability.


## Boltzmann Samplers

The size of an object generated by $\Gamma C(\lambda)$ is a random variable denoted by $U$ with the probability distribution

$$
\mathbb{P}_{\lambda}(U=n)=\frac{c_{n} \lambda^{n}}{C(\lambda)}
$$



Figure: Distribution for integer partitions ( $n=500, \lambda=0.90$ ).

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Figure: Distribution for integer partitions ( $n=500, \lambda=0.91$ ).

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Figure: Distribution for integer partitions ( $n=500, \lambda=0.92$ ).

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Figure: Distribution for integer partitions ( $n=500, \lambda=0.94$ ).

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Figure: Distribution for integer partitions ( $n=500, \lambda=0.95$ ).

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Figure: Distribution for integer partitions ( $n=500, \lambda \in[0.90,0.95]$ ).

## Sampling Algorithm for Weighted Partitions

Def. Let $\left(b_{k}\right)_{k=1}^{\infty}$ be a sequence such that $b_{k}=p(k)$ for some polynomial $p(x)=a_{0}+a_{1} x+\cdots+a_{r} x^{r} \in \mathbb{R}[x]$ with $\operatorname{deg}(p)=r$.

- $\binom{k+2}{2}=\frac{1}{2}(k+1)(k+2)$

Algorithm 1 Sampling algorithm for weighted partitions.
1: procedure RandomWeightedPartition $(n)$
2: $\quad \lambda_{n} \leftarrow$ Solution to $\sum_{k=1}^{n} k b_{k} \lambda^{k} /\left(1-\lambda^{k}\right)=n \quad \triangleright$ Tuning
3: repeat
$\triangleright$ Rejection sampling
$\begin{array}{ll}\text { 4: } & \quad \gamma \leftarrow \Gamma C_{n}\left(\lambda_{n}\right) \\ \text { 5: } & \text { until }|\gamma|=n \\ \text { 6: } & \text { return } \gamma\end{array}$

Theorem. Algorithm 1 runs in expected $O\left(n^{r+1}((r+2) n)^{3 / 4}\right)$ time and uses $O(n)$ space.

## Sampling Algorithm for Weighted Partitions

Observation.

$$
\mathbb{P}_{\lambda}(\gamma)=\frac{\lambda^{|\gamma|}}{\prod_{k=1}^{n}\left(1-\lambda^{k}\right)^{-b_{k}}}=\prod_{k=1}^{n} \prod_{j=1}^{b_{k}} \frac{\lambda^{\left(\# \text { of columns of type } b_{k, j}\right) k}}{\left(1-\lambda^{k}\right)^{-1}}
$$

Algorithm 2 Boltzmann sampler for weighted partitions.
1: procedure $\Gamma C_{n}(\lambda)$
2: $\quad \gamma \leftarrow$ Empty associative array
3: $\quad$ for $k=1$ to $n$ do
4: $\quad$ for $j=1$ to $b_{k}$ do
5: $\quad m \leftarrow$ Geometric $\left(1-\lambda^{k}\right)$
6: $\quad$ if $m \geq 1$ then
$7:$

$$
\gamma[(k, j)] \leftarrow m
$$

8: $\quad$ return $\gamma$

- The PDF of Geometric $\left(1-\lambda^{k}\right)$ is $\mathbb{P}_{1-\lambda^{k}}(m)=\lambda^{m k}\left(1-\lambda^{k}\right)$.


## Tuning the Boltzmann Samplers

The random variable for the size of an object produced is

$$
U_{n}=\sum_{k=1}^{n} \sum_{j=1}^{b_{k}} k Y_{k, j}
$$

where $Y_{k, j} \sim \operatorname{Geometric}\left(1-\lambda^{k}\right)$.

Lemma. We have

$$
\mathbb{E}_{\lambda}\left[U_{n}\right]=\sum_{k=1}^{n} k b_{k}\left(\frac{\lambda^{k}}{1-\lambda^{k}}\right)
$$

- Follows from the linearity of expectation and the mean of geometric random variable
- Strictly increasing so binary search to solve $\mathbb{E}_{\lambda}\left[U_{n}\right]=n$


## Bounding the Rejection Rates

Lemma. The Dirichlet generating series for weighted partitions is

$$
D(s)=\sum_{k=1}^{\infty} \frac{b_{k}}{k^{s}}=\sum_{k=0}^{r} a_{k} \zeta(s-k)
$$

and has at most $r+1$ simple poles on the positive real axis at positions $\rho_{k}=k+1$ with residue $A_{k}=a_{k}$ if and only if $a_{k} \neq 0$.

Proof. The Riemann zeta function converges uniformly and is analytic on $\mathbb{C} \backslash\{1\}$, so

$$
D(s)=\sum_{k=1}^{\infty} \frac{b_{k}}{k^{s}}=\sum_{k=1}^{\infty} \frac{a_{0}+a_{1} k+\cdots+a_{r} k^{r}}{k^{s}}=\sum_{k=0}^{r} a_{k} \zeta(s-k) .
$$

Since $\zeta(s)=\sum_{k=1}^{\infty} 1 / k^{s}$ has a simple pole at $s=1$ with residue 1 , the claim about the poles of $D(s)$ follows.

## Bounding the Rejection Rates

Theorem [Granovsky and Stark, 2012]. For $n$ sufficiently large,

$$
\mathbb{P}_{\lambda_{n}}\left(U_{n}=n\right) \geq \frac{1}{10}\left(\frac{A_{r}}{2} \Gamma\left(\rho_{r}+2\right) \zeta\left(\rho_{r}+1\right)\right)^{\frac{1}{2(r+2)}}((r+2) n)^{-3 / 4}
$$

Lemma. For any positive integer sequence of degree $r$,

$$
\frac{A_{r}}{2} \Gamma\left(\rho_{r}+2\right) \zeta\left(\rho_{r}+1\right) \geq 1
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Proof. Let

$$
p(k)=\sum_{j=0}^{r} \Delta^{j} p(0)\binom{k}{j}
$$

[Stanley, EC1, Cor 1.9.3] $\Longrightarrow \Delta^{r} p(0) \in \mathbb{Z} \quad \Longrightarrow \quad a_{r}=A_{r} \geq \frac{1}{r!}$. Since $\rho_{r}=r+1$,

$$
\frac{A_{r}}{2} \Gamma\left(\rho_{r}+2\right) \zeta\left(\rho_{r}+1\right) \geq \frac{1}{2 r!}(r+2)!\geq 1
$$

## Boltzmann Sampler for Bose-Einstein Condensation

Lemma. If $b_{k}=\binom{k+d-1}{d-1}$, for $d \geq 1$, then

$$
C(z)=\prod_{k=1}^{\infty}\left(1-z^{k}\right)^{-\binom{k+d-1}{d-1}}=\prod_{k=1}^{\infty} \exp \left(\frac{1}{k}\left[\left(1-z^{k}\right)^{-d}-1\right]\right)
$$

Theorem. We can uniformly sample objects of size $n$ in $\mathcal{C}$ in expected $O\left(n((d+1) n)^{3 / 4}\right)$ time while using $O(n)$ space.

- Runtime reduced from $O\left(n^{3.75}\right)$ to $O\left(n^{1.75}\right)$
- Columns are sampled from the zero-truncated negative binomial distribution
- Geometric random variable is a weighted sum of independent Poisson random variables


## References

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