# Planar graphs with girth at least 5 are (3, 4)-colorable

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• On proper coloring of planar graphs, a famous example is the Four Color Theorem.

• We may relax the requirement by allowing some edges in each color class.

A graph G is called (d<sub>1</sub>, d<sub>2</sub>,..., d<sub>r</sub>)-colorable, if its vertex set can be partitioned into r nonempty subsets so that the subgraph induced by the *i*th part has maximum degree at most d<sub>i</sub> for each i ∈ {1,...,r}, where d<sub>i</sub>s are non-negative integers.

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- Improper colorings have then been considered for planar graphs with large girth or graphs with low maximum average degree. (See Montassier and Ochem, Near-colorings: non-colorable graphs and NP-completeness, the electronic journal of combinatorics 22(1) (2015), #P1.57)

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- The Four Color Theorem says that every planar graph is (0,0,0,0)-colorable.
- In 1986, Cowen, Cowen, and Woodall proved that planar graphs are (2, 2, 2)-colorable. In 1999, Eaton and Hull, Škrekovski, separately, proved that this is sharp by exhibiting non-(1, k, k)-colorable planar graphs for each k. Thus, the problem is completely solved when r ≥ 3.

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- There are non-(3, 1)-colorable planar graphs in  $\mathcal{G}_5$ . (Montassier and Ochem, 2015)

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- Planar graphs in  $\mathcal{G}_5$  are (3, 5)-colorable. (Choi and Raspaud 2015)
- Planar graphs in  $\mathcal{G}_5$  are (4,4)-colorable. (Havet and Sereni 2006)

# A summary on $(d_1, d_2)$ -coloring

| girth | ( <i>k</i> ,0) | (k, 1)  | (k,2) | ( <i>k</i> ,3) | ( <i>k</i> , 4) |
|-------|----------------|---------|-------|----------------|-----------------|
| 3,4   | х              | х       | х     | х              | х               |
| 5     | Х              | (10, 1) | (6,2) | (5,3)          | (4,4)           |
| 6     | Х              | (4,1)   | (2,2) |                |                 |
| 7     | (4,0)          | (1, 1)  |       |                |                 |
| 8     | (2,0)          |         |       |                |                 |
| 11    | (1,0)          |         |       |                |                 |

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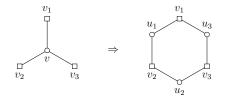
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- Claim. G must be connected and there are no 1-vertices in G.
- Lemma 2 There is no 3-vertex in G.



• Let the initial charge of each element  $x \in V \cup F$  be  $\mu(x) = d(x) - 4$ . Then by Euler formula,

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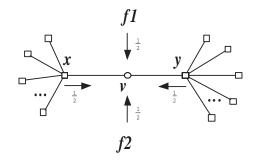
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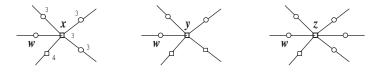
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- Clearly, by Lemma 2, each face and each vertex has a non-negative initial charge except 2-vertices.

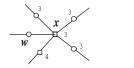
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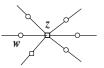
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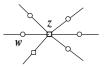


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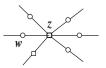
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$$\mu'(z) = 6 - 4 - \frac{5}{2} = -\frac{1}{2}$$
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## Proof+

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- Considering these three special vertices and some special faces, we design the discharging rules.
- By the discharging rules, there is µ<sup>\*</sup>(x) ≥ 0 for each x ∈ V ∪ F. So we have

$$\sum_{x\in V\cup F}\mu^*(x)\geq 0,$$

a contradiction.

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- **Problem 3.** What is the minimum *d* where graphs with girth 5 are (3, *d*)-colorable in {2,3,4}?

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Thank you for your attention!